

# Automorphism groups and adversarial vertex deletions

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## Abstract

Any finite group can be encoded as the automorphism group of an unlabeled simple graph. Recently Hartke, Kolb, Nishikawa, and Stolee (2010) demonstrated a construction that allows any ordered pair of finite groups to be represented as the automorphism group of a graph and a vertex-deleted subgraph. In this note, we describe a generalized scenario as a game between a player and an adversary: an adversary provides a list of finite groups and a number of rounds. The player constructs a graph with automorphism group isomorphic to the first group. In the following rounds, the adversary selects a group and the player deletes a vertex such that the automorphism group of the corresponding vertex-deleted subgraph is isomorphic to the selected group. We provide a construction that allows the player to appropriately respond to any sequence of challenges from the adversary.

## 1 Introduction

Automorphisms of graphs are incredibly unstable. The slightest perturbation of the graph can greatly change the automorphism group. In this note, we show there exist graphs whose automorphism groups can change dramatically under certain sequences of vertex deletions. We consider undirected, unlabeled, and simple graphs, denoted  $F$ ,  $G$ , or  $H$ , and finite groups, denoted  $\Gamma$ . The automorphism group of a graph  $G$  is denoted  $\text{Aut}(G)$ .

Frucht [3] proved that graphs have the ability to encode the structure of any finite group.

**Theorem 1** (Frucht [3]). *Let  $\Gamma$  be a finite group. There exists a graph  $G$  with  $\text{Aut}(G) \cong \Gamma$ .*

Hartke, Kolb, Nishikawa, and Stolee [4] proved that any ordered pair of finite groups can be represented by a graph and a vertex-deleted subgraph. Their work was motivated by consequences to the Reconstruction Conjecture (see Bondy [2]) and isomorph-free generation (see McKay [5]).

**Theorem 2** (Hartke, Kolb, Nishikawa, Stolee [4]). *Let  $\Gamma_0$  and  $\Gamma_1$  be finite groups. There exists a graph  $G$  and a vertex  $v \in V(G)$  such that  $\text{Aut}(G) \cong \Gamma_0$  and  $\text{Aut}(G - v) \cong \Gamma_1$ .*

There are two natural extensions of this process to a sequence  $\Gamma_0, \Gamma_1, \dots, \Gamma_k$  of finite groups using two types of vertex deletions: single deletions or iterated deletions.

**Question.** *Let  $\Gamma_0, \Gamma_1, \dots, \Gamma_k$  be finite groups. Does there exist a graph  $G$  with vertices  $v_1, \dots, v_k \in V(G)$  such that  $\text{Aut}(G) \cong \Gamma_0$  and for all  $i \in \{1, \dots, k\}$ ,*

1. *(Single Deletions)  $\text{Aut}(G - v_i) \cong \Gamma_i$ ?*
2. *(Iterated Deletions)  $\text{Aut}(G - v_1 - \dots - v_i) \cong \Gamma_i$ ?*

In fact, both of these types of deletions can be combined in an even more general situation, posed as the *vertex deletion game* between a player and an adversary:

### The Vertex Deletion Game

**Round 0:**

*Adversary:* Selects finite groups  $\Gamma_0, \Gamma_1, \dots, \Gamma_k$ , and a number  $\ell \geq 1$ .

*Player:* Constructs a graph  $G_0$  with  $\text{Aut}(G_0) \cong \Gamma_0$ .

**Round  $j$ :**  $(1 \leq j \leq \ell)$

*Adversary:* Selects a group  $\Gamma_{i_j} \in \{\Gamma_1, \dots, \Gamma_k\}$ .

*Player:* Selects a vertex  $v_j \in V(G_{j-1})$ , defines  $G_j = G_{j-1} - v_j$ , and asserts  $\text{Aut}(G_j) \cong \Gamma_{i_j}$ .

Note that this game generalizes both single deletions (play the game with  $\ell = 1$ ) and iterated deletions (play the game with  $\ell = k$ , and the adversary selects  $\Gamma_{i_j} = \Gamma_j$  for all  $j \in \{1, \dots, k\}$ ). By carefully constructing  $G_0$ , the player can survive  $\ell$  rounds against the adversary.

**Theorem 3** (Adversarial Iterated Deletions). *Suppose the adversary selects  $\Gamma_0, \Gamma_1, \dots, \Gamma_k$  as finite groups and integer  $\ell \geq 1$  in Round 0. The player can construct a graph  $G_0$  with  $\text{Aut}(G_0) \cong \Gamma_0$  so that the assertions  $\text{Aut}(G_j) \cong \Gamma_{i_j}$  hold for all  $\ell$  remaining rounds.*

Instead of using the vertex deletion game, there is an equivalent statement of the previous theorem using a sequence of alternating quantifiers.

**Theorem 4** (Adversarial Iterated Deletions; alternate form). *For all numbers  $k, \ell \geq 1$  and finite groups  $\Gamma_0, \Gamma_1, \dots, \Gamma_k$ , there exists a graph  $G_0$  such that  $\text{Aut}(G_0) \cong \Gamma_0$  and*

$$\forall i_1 \exists v_1 \forall i_2 \exists v_2 \cdots \forall i_\ell \exists v_\ell \forall j, \text{Aut}(G_0 - v_1 - \cdots - v_j) \cong \Gamma_{i_j},$$

where the domain of  $j$  is  $\{1, \dots, \ell\}$ , the domain of each  $i_j$  is  $\{1, \dots, k\}$ , and the domain of each  $v_j$  is  $V(G_0) \setminus \{v_1, \dots, v_{j-1}\}$ .

A group is *trivial* if it consists only of the identity element. For a graph  $G$  and vertex  $v \in V(G)$ , the *stabilizer* of  $v$  in  $G$ , denoted  $\text{Stab}_G(v)$ , is the subgroup of  $\text{Aut}(G)$  given by permutations  $\tau$  where  $\tau(v) = v$ .

## 2 Results

Our starting point is the following lemma from [4].

**Lemma 5** (Hartke, Kolb, Nishikawa, Stolee [4, Lemma 2.2]). *For any finite group  $\Gamma$ , there is a connected graph  $G$  and a vertex  $v \in V(G)$  where  $\text{Aut}(G) \cong \Gamma$  and  $\text{Stab}_G(v)$  is trivial.*

The construction for Lemma 5 from [4] uses the Cayley graph for  $\Gamma$  and replaces labeled, directed edges with undirected gadgets. This results with  $G$  having maximum degree  $2|\Gamma|$  and  $|V(G)| \geq 5|\Gamma|$ . This specific construction contains no vertices in  $G$  of degree  $|V(G)| - 1$ , which we will use in the following lemma.

We now describe a gadget which will be used to build the full construction for Theorem 3.

**Lemma 6.** *Let  $\Gamma$  be a finite group. There exists a graph  $H$  and two vertices  $x, y \in V(H)$  so that  $\text{Aut}(H)$  is trivial,  $H - x$  is connected,  $\text{Aut}(H - x) \cong \Gamma$ , and  $\text{Stab}_{H-x}(y)$  is trivial.*

*Proof.* By Lemma 5, there exists a connected graph  $G$  and a vertex  $y \in V(G)$  so that  $\text{Aut}(G) \cong \Gamma$  and  $\text{Stab}_G(y)$  is trivial. Let  $n = |V(G)|$  and order the vertices of  $G$  as  $V(G) = \{v_1, \dots, v_n\}$  and  $v_1 = y$ .

Let  $H$  be a graph with vertex set  $V(H) = \{v_1, \dots, v_n\} \cup \{u_1, \dots, u_n\} \cup \{x, z, w\}$ . The graph  $H$  has an edge  $v_i v_j$  if and only if that edge is present in  $G$ . For every  $j \in \{1, \dots, n\}$ , the pair  $u_j v_j$  is an edge. The vertex  $z$  is adjacent to all vertices  $v_j$  for  $j \in \{1, \dots, n\}$ . The vertex  $x$  is adjacent to  $z$ , all vertices  $v_j$  for  $j \in \{1, \dots, n\}$  and adjacent to the vertices  $u_i$  for  $i \in \{2, \dots, n\}$ . Finally, the vertex  $w$  is adjacent only to  $x$  and  $z$ .

The only vertex of degree 1 in  $H$  is  $u_1$ , so every automorphism of  $H$  stabilizes  $u_1$  and thus also stabilizes  $v_1$ . All vertices  $v_1, \dots, v_n$  have degree at least three and degree at most  $n + 1$ . The vertex  $x$  is the only vertex of degree  $2n + 1$ , so every automorphism of  $H$  stabilizes  $x$ . The vertex  $z$  is the only vertex of degree  $n + 2$ , so every automorphism of  $H$  stabilizes  $z$  and hence also stabilizes  $w$ . Other than  $w$ , the vertices  $u_2, \dots, u_n$  are the only vertices of degree 2, so every automorphism

set-wise stabilizes  $\{u_2, \dots, u_n\}$ . This implies that every automorphism of  $H$  set-wise stabilizes  $\{v_1, \dots, v_n\}$  and hence restricting an automorphism of  $H$  to  $V(G)$  induces an automorphism of  $G$ . However, every automorphism point-wise stabilizes  $v_1$ . Since  $v_1 = y$  and  $\text{Stab}_G(y)$  is trivial, every automorphism of  $H$  must point-wise stabilize  $V(G)$ . Thus the vertices  $u_1, \dots, u_n$  are also point-wise stabilized and the automorphism group of  $H$  is trivial.

Now consider  $H' = H - x$ . The vertex  $z$  is the only vertex of degree  $n + 1$ , so every automorphism of  $H'$  stabilizes  $z$  and  $w$ . Other than  $w$ , the only other vertices of degree 1 are  $u_1, \dots, u_n$ , and since  $w$  is the only one of these that is adjacent to the stabilized  $z$ , every automorphism of  $H'$  set-wise stabilizes  $w$  and therefore set-wise stabilizes  $\{u_1, \dots, u_n\}$ . Since the vertices  $u_1, \dots, u_n$  are adjacent only to vertices in  $V(G)$ , every automorphism of  $H'$  set-wise stabilizes  $V(G)$ . Hence, every automorphism of  $H'$  restricted to  $V(G)$  is an automorphism of  $G$ . Observe that every automorphism  $\sigma \in \text{Aut}(G)$  extends to an automorphism of  $H$  by assigning  $\sigma(u_i) = u_j$  whenever  $\sigma(v_i) = v_j$ . Thus,  $\text{Aut}(H - x) \cong \text{Aut}(G) \cong \Gamma$ .

Since the automorphisms of  $H - x$  correspond directly to automorphisms of  $G$ , observe that  $\text{Stab}_{H-x}(y)$  is trivial.  $\square$

We are now sufficiently prepared to prove the main theorem. The gadget from Lemma 6 has two purposes:

1. “Reveal” symmetry: When  $x$  is deleted, the automorphism group  $\Gamma$  is revealed.
2. “Remove” symmetry: When  $y$  is stabilized within  $H - x$ , all non-trivial automorphisms of  $H - x$  are removed.

Our construction for the graph  $G_0$  carefully places many copies of this gadget in such a way that the player has access to a “revealing” vertex ( $x$ ) that simultaneously stabilizes the “removing” vertex ( $y$ ) in the previous gadget. Therefore, we have a sequence of deletions which remove all previous symmetry and reveal only the requested symmetry.

*Proof of Theorem 3.* Note that the case  $k = \ell = 1$  holds by Theorem 2. Clearly, we may assume that the groups  $\Gamma_1, \dots, \Gamma_k$  are distinct with respect to isomorphism.

By Lemma 6, for every  $i \in \{0, 1, \dots, k\}$  there is a graph  $H_i$  with vertices  $x_i, y_i \in V(H_i)$  such that  $\text{Aut}(H_i)$  is trivial,  $\text{Aut}(H_i - x_i) \cong \Gamma_i$ , and  $\text{Stab}_{H_i-x_i}(y_i)$  is trivial. For all  $i \in \{0, \dots, k\}$ , let  $O_i$  be the orbit of  $y_i$  in  $H_i - x_i$ . Since the groups  $\Gamma_1, \dots, \Gamma_k$  are pairwise non-isomorphic, then by the construction of Lemma 6 the graphs  $H_0, H_1, \dots, H_k$  and  $H_0 - x_0, \dots, H_k - x_k$  are all pairwise non-isomorphic. Also by the construction of Lemma 6, no graph  $H_i$  or  $H_i - x_i$  has a dominating vertex, and  $x_i$  is a maximum-degree vertex of  $H_i$ . Hence for all  $i \neq j$ , if an isomorphism  $\phi_{ij} : H_i \rightarrow H_j$  exists, then  $\phi_{ij}$  sends  $x_i$  to  $x_j$  and hence  $\phi_{ij}$  induces an isomorphism  $\psi_{ij} : H_i - x_i \rightarrow H_j - x_j$ . However, if  $H_i - x_i \cong H_j - x_j$  then  $\Gamma_i \cong \Gamma_j$ , which contradicts our assumption that the groups are pairwise nonisomorphic. Therefore,  $H_i \not\cong H_j$  and  $H_i - x_i \not\cong H_j - x_j$  for all  $i \neq j$ .

We construct the graph  $G_0$  by building graphs  $F_0, F_1, \dots, F_\ell$  iteratively. Let  $F_0$  be the graph given by taking  $H_0 - x_0$  and adding vertices  $a_0, b_0$  where  $a_0$  is adjacent to all vertices in  $H_0 - x_0$  and  $b_0$  is adjacent to only  $a_0$ . Let  $U_0 = O_0$ .

For all  $j \in \{1, \dots, \ell\}$ , we will build  $F_j$  by adding vertices and edges to  $F_{j-1}$ . During the process,  $F_{j-1}$  will remain an induced subgraph of  $F_j$ . For all vertices  $v \in U_{j-1}$  and  $i \in \{1, \dots, k\}$ , add a copy  $H_i^{(j,v)}$  of  $H_i$  to  $F_{j-1}$  and add edges from  $v$  to each vertex of  $H_i^{(j,v)}$ . Let  $x_i^{(j,v)}$  and  $y_i^{(j,v)}$  denote the copies of  $x_i$  and  $y_i$  in  $H_i^{(j,v)}$ . Let  $O_i^{(j,v)}$  be the copy of  $O_i$  within  $H_i^{(j,v)}$  and define  $U_j = \cup_{v \in U_{j-1}} \cup_{i=1}^k O_i^{(j,v)}$ . Add vertices  $a_j, b_j$  where  $a_j$  is adjacent to all vertices in  $V(F_j) \setminus V(F_{j-1})$  and the vertices  $a_{j-1}$  and  $b_j$ . Figure 1 shows a visualization of this construction.

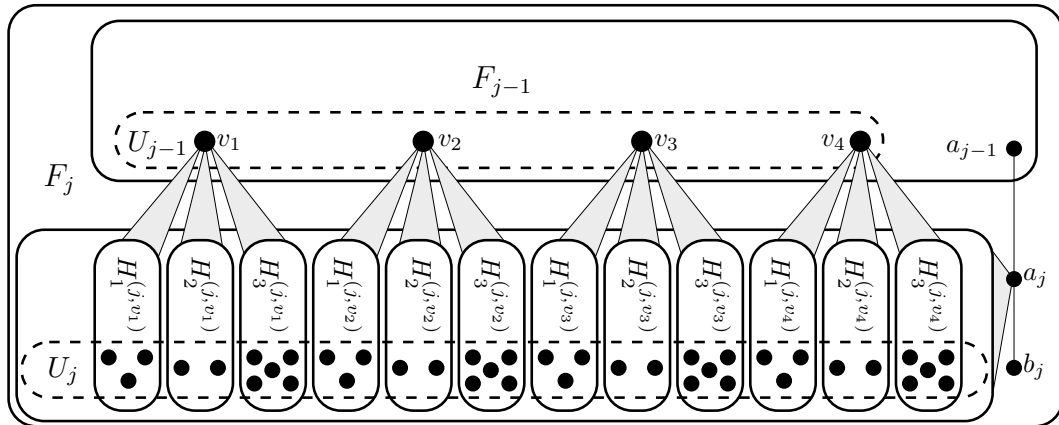


Figure 1: An example of the construction of  $F_j$  from  $F_{j-1}$  where  $k = 3$ .

Let  $G_0$  be  $F_\ell$ . Observe that the vertices  $a_0, \dots, a_\ell$  induce a path, and the vertices  $b_0, \dots, b_\ell$  all have degree 1. The vertices  $b_0, \dots, b_\ell$  are the only vertices of degree 1, so all automorphisms of  $G_0$  set-wise stabilize  $\{b_0, \dots, b_\ell\}$  and hence set-wise stabilize  $\{a_0, \dots, a_\ell\}$ . Since all vertices in  $\{a_0, \dots, a_\ell\}$  have distinct degrees, these vertices are point-wise stabilized by all automorphisms of  $G_0$ . Therefore, every set  $V(F_j) \setminus V(F_{j-1})$  is set-wise stabilized by every automorphism of  $G_0$ . In particular, any automorphisms of  $G_0$  must set-wise stabilize the set  $V(F_0) - \{a_0, b_0\}$  which induces a copy of  $H_0 - x_0$ .

It remains to show that  $G_0$  satisfies the conditions of Theorem 3 by providing a strategy for the player to respond to the adversary’s challenges. Informally, in the  $j$ th round the player will delete a vertex from the  $j$ th layer (i.e.  $V(F_j) \setminus V(F_{j-1})$ ), and this vertex will depend on the  $j$ th group,  $\Gamma_{i_j}$ , and the previous vertex-deletions. The previous vertex-deletion removed a copy of the vertex  $x_{i_{j-1}}$  from a copy of  $H_{i_{j-1}}$  (or  $j = 1$ ,  $i_0 = 0$ , and  $x_0$  was never included in  $G_0$ ). To “remove” the symmetry found in this copy of  $H_{i_{j-1}}$ , we aim to stabilize its copy of  $y_{i_{j-1}}$ . We delete the vertex  $x_{i_j}$  from the copy of  $H_{i_j}$  in the neighborhood of this copy of  $y_{i_{j-1}}$ , which distinguishes it from all other vertices in  $U_{j-1}$  and hence the symmetry in  $H_{i_{j-1}} - x_{i_{j-1}}$

is no longer available. Instead, we have removed  $x_{i_j}$  from a copy of  $H_{i_j}$ , revealing  $\text{Aut}(H_{i_j} - x_{i_j}) \cong \Gamma_{i_j}$ . Thus, the automorphisms allowed within  $F_j$  are exactly those in this copy of  $H_{i_j} - x_{i_j}$ , and all vertices in  $F_\ell \setminus F_j$  have their motion determined by the action in  $F_j$ .

Back to the formal proof, we first show that we can localize our study of the automorphisms of  $G_0 - X$  for certain sets of vertices  $X$ .

**Claim 7.** *Fix  $j \in \{0, \dots, \ell\}$  and  $X = \{v_1, \dots, v_j\}$  where  $v_{j'} \in V(F_{j'}) \setminus (V(F_{j'-1}) \cup U_{j'} \cup \{a_{j'}, b_{j'}\})$  for all  $j' \in \{1, \dots, j\}$ . Then  $\text{Aut}(G_0 - X) \cong \text{Aut}(F_j - X)$ .*

*Proof of Claim 7:* Observe that the vertices  $b_0, \dots, b_\ell$  remain the only vertices in  $G_0 - X$  of degree 1, and the vertices  $a_0, \dots, a_\ell$  continue to have distinct degrees. Thus, the vertices  $a_0, \dots, a_\ell$  are point-wise stabilized by  $\text{Aut}(G_0 - X)$  and hence the sets  $V(F_i) \setminus V(F_{i-1})$  are set-wise stabilized by  $\text{Aut}(G_0 - X)$ . Specifically, the sets  $V(F_{\ell'+1}) \setminus V(F_{\ell'})$  are set-wise stabilized by  $\text{Aut}(G_0 - X)$  for all  $\ell' \in \{j, \dots, \ell - 1\}$ . This implies that every automorphism in  $\text{Aut}(F_{\ell'+1} - X)$  is also an automorphism of  $\text{Aut}(F_{\ell'} - X)$  when restricted to  $V(F_{\ell'} - X)$ .

We will show that this map from  $\text{Aut}(F_{\ell'+1} - X)$  to  $\text{Aut}(F_{\ell'} - X)$  is a bijection for all  $\ell' \in \{j, \dots, \ell - 1\}$ , implying there is natural bijection between  $\text{Aut}(F_\ell - X)$  and  $\text{Aut}(F_j - X)$ . Every vertex  $u \in V(F_{\ell'+1} - X) \setminus V(F_{\ell'})$  is contained in  $H_i^{(\ell'+1, v)}$  for some vertex  $v \in U_{\ell'}$  and  $i \in \{1, \dots, k\}$ . Since  $X \subset V(F_j)$  and  $\ell \geq j$ , it follows that  $V(H_i^{(\ell'+1, v)}) \cap X = \emptyset$ . Therefore, the subgraph  $H_i^{(\ell'+1, v)}$  is a copy of  $H_i$  in  $G_0 - X$ , and has no non-trivial automorphisms by Lemma 6. Therefore, for every automorphism  $\sigma$  of  $F_{\ell'} - X$ , there is exactly one isomorphism of  $F_{\ell'+1} - X$  that extends  $\sigma$  and maps  $V(H_i^{(\ell'+1, v)})$  to  $V(H_i^{(\ell'+1, \sigma(v))})$ . Hence, the action of an automorphism on each vertex  $u \in V(F_{\ell'+1} - X) \setminus V(F_{\ell'})$  is determined exactly by the action of the automorphism on the vertices within  $V(F_{\ell'} - X)$ . Hence, the restriction map from  $\text{Aut}(F_{\ell'+1} - X)$  to  $\text{Aut}(F_{\ell'} - X)$  is a bijection, proving the claim.  $\square$

When  $X = \emptyset$ , the automorphism group of the subgraph  $F_0$  determines the automorphism group of  $G_0 - X$ . Since  $F_0 - \{a_0, b_0\} \cong H_0 - x_0$ , we have  $\text{Aut}(G_0) \cong \Gamma_0$ .

For a list  $\Gamma_{i_1}, \dots, \Gamma_{i_\ell}$  of groups selected from  $\{\Gamma_1, \dots, \Gamma_k\}$ , define the vertices  $v_1, \dots, v_\ell$  and  $u_0, u_1, \dots, u_\ell$  where  $u_0 = y_0$  and for  $j \in \{1, \dots, \ell\}$ ,

$$v_j = x_{i_j}^{(j, u_{j-1})}, \quad u_j = y_{i_j}^{(j, u_{j-1})}.$$

Observe that the definition of  $v_j$  and  $u_j$  depends only on  $u_{j-1}$  and  $\Gamma_{i_j}$ , so this definition does not violate any of the quantifiers in the statement of Theorem 4. Thus, the vertices  $v_1, \dots, v_\ell$  are valid selections of vertex-deletions for the player in response to the adversary selecting  $\Gamma_{i_1}, \dots, \Gamma_{i_\ell}$  in order.

By induction on  $j$ , we verify that  $\text{Aut}(G_0 - v_1 - \dots - v_j) \cong \Gamma_{i_j}$ . We will require the stronger induction hypothesis that all automorphisms of  $F_j - v_1 - \dots - v_j$  point-wise stabilize all vertices except those in  $H_{i_j}^{(j, u_{j-1})} - v_j$ .

Let  $j \in \{1, \dots, \ell\}$ . By Claim 7,  $\text{Aut}(G_0 - v_1 - \dots - v_j) \cong \text{Aut}(F_j - v_1 - \dots - v_j)$ . Since  $b_0, \dots, b_j$  are the only vertices of degree 1, and they are only adjacent to  $a_0, \dots, a_j$  (which have different degrees), the vertices  $a_0, \dots, a_j$  and  $b_0, \dots, b_j$  are

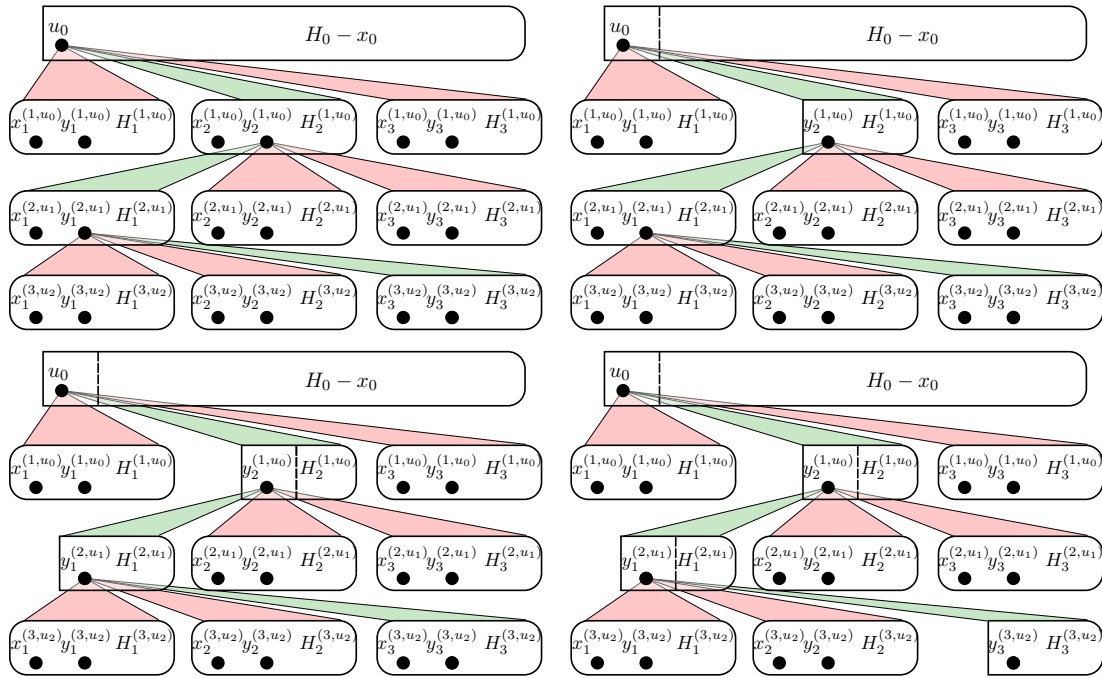


Figure 2: An example sequence of deletions with  $k = 3$  where the adversary selects  $\Gamma_2, \Gamma_1, \Gamma_3$ , we define  $v_1 = x_2^{(1,u_0)}, u_1 = y_2^{(1,v_1)}, v_2 = x_1^{(2,u_1)}, u_2 = y_1^{(2,u_1)}, v_3 = x_3^{(3,u_2)}, u_3 = y_3^{(3,u_2)}$ , and the player deletes  $v_1, v_2, v_3$ .

point-wise stabilized by  $\text{Aut}(F_j - v_1 - \dots - v_j)$ . Thus,  $V(F_{j-1})$  is set-wise stabilized by  $\text{Aut}(F_j - v_1 - \dots - v_j)$ . By induction (or that  $F_0 - \{a_0, b_0\} \cong H_0 - x_0$  in the case  $j = 1$ ),  $\text{Aut}(F_{j-1} - v_1 - \dots - v_{j-1}) \cong \Gamma_{i_{j-1}}$  and all vertices in  $F_{j-1}$  are point-wise stabilized by  $\text{Aut}(F_{j-1} - v_1 - \dots - v_{j-1})$  except those in  $H_{i_{j-1}}^{(j-1, u_{j-2})} - v_{j-1}$  (for the case  $j = 1$ , use  $H_0 - x_0$  instead of  $H_{i_{j-1}}^{(j-1, u_{j-2})} - v_{j-1}$ ). Observe that  $u_{j-1}$  is the copy of  $y_{i_{j-1}}$  in  $H_{i_{j-1}}^{(j-1, u_{j-2})}$ . Since deleting  $v_j$  from  $F_j - v_1 - \dots - v_j$  creates a copy of  $H_{i_j} - x_{i_j}$  in the neighborhood of  $u_{j-1}$ , the vertex  $u_{j-1}$  is distinguished from the other vertices in  $U_{j-1}$ . Thus,  $u_{j-1}$  is stabilized by all automorphisms in  $\text{Aut}(F_j - v_1 - \dots - v_j)$ . This implies that the automorphisms in  $\text{Aut}(F_j - v_1 - \dots - v_j)$  point-wise stabilize all vertices in  $F_{j-1}$ . Finally, all vertices in  $V(F_j) \setminus V(F_{j-1})$  are either contained in  $H_{i_j}^{(j, u_{j-1})} - v_j$  (in which case the automorphisms are given by  $\text{Aut}(H_{i_j} - x_{i_j})$ ) or are contained in a copy of  $H_i$  for some  $i \in \{1, \dots, k\}$  and  $H_i$  has no nontrivial automorphisms. Thus, all vertices of  $F_j - v_1 - \dots - v_j$  are point-wise stabilized except those in  $H_{i_j}^{(j, u_{j-1})} - v_j$ . Finally,

$$\text{Aut}(G_0 - v_1 - \dots - v_j) \cong \text{Aut}(F_j - v_1 - \dots - v_j) \cong \text{Aut}(H_{i_j}^{(j, u_{j-1})} - v_j) \cong \Gamma_{i_j}. \quad \square$$

The construction given in the above proof requires a large number of vertices and vertices of high degree. The gadget given by Lemma 6 can be built using  $O(|\Gamma| \log_2^2 |\Gamma| \log_2 \log_2 |\Gamma|)$  vertices: the construction of Lemma 5 from [4] has order  $O(|\Gamma|^4)$ , but can be replaced by a construction of Sabidussi [6] with

$O(|\Gamma| \log_2 |\Gamma| \log_2 \log_2 |\Gamma|)$  vertices, then carefully applying the construction of Lemma 6 to Sabidussi's construction increases the number of vertices by a multiplicative factor of  $O(\log_2 |\Gamma|)$ . However, Babai [1] proved that for every finite group  $\Gamma$  there is a graph  $G$  with  $\text{Aut}(G) \cong \Gamma$  and  $|V(G)| \leq 3|\Gamma|$ . Can graphs with  $O(|\Gamma|)$  vertices be used to satisfy Lemma 6? Also, the constructions used here contain vertices of high degree. Does there exist a constant  $D$  so that Theorem 3 is satisfied with the maximum degree of  $G_0$  at most  $D$ ?

## Acknowledgements

The author thanks the anonymous referees who provided several comments that greatly improved this paper, especially an improvement to the construction in Lemma 6. Also, thanks to Michael Ferrara for convincing the author to write the paper.

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(Received 26 June 2013; revised 23 Apr 2014, 14 May 2014)